Bounds on some van der Waerden numbers

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Abstract

For positive integers *s* and $k_1; k_2; ...; k_s$, the van der Waerden number $w(k_1; k_2; ...; k_s; s)$ is the minimum integer *n* such that for every *s*-coloring of set $f_1; 2; ...; ng$, with colors 1; 2; ...; s, there is a k_i -term arithmetic progression of color *i* for some *i*. We give an asymptotic lower bound for w(k;m;2) for fixed *m*. We include a table of values of w(k;3;2) that are very close to this lower bound for m = 3. We also give a lower bound for w(k;k;...;k;s) that slightly improves previously-known bounds. Upper bounds for w(k;4;2) and w(4;4;...;4;s) are also provided.

1 Introduction

Two fundamental theorems in combinatorics are van der Waerden's theorem [18] and Ramsey's theorem [16]. The theorem of van der Waerden says that for all positive integers *s* and k_1 ; k_2 ; ...; k_s , there exists a least positive integer $n = w(k_1; k_2; ...; k_s; s)$ such that whenever $[1; n] = f_1; 2; ...; ng$ is *s*-colored (i.e., partitioned into *s* sets), there is a k_i -term arithmetic progression with color *i* for some *i*, 1 = i = s.

Similarly, Ramsey's theorem has an associated "threshold" function $R(k_1; k_2; ...; k_s; s)$ (which we will not define here). The order of magnitude of R(k;3;2) is known to be $\frac{k^2}{\log k}$ [11], while the best known upper bound on R(k;m;2) is fairly close to the best known lower bound. In contrast, the order of magnitude of w(k;3;2) is not known, and the best known lower and upper bounds on w(k;k]?

of previously known values w(3;3;2) = 9, w(4;4;2) = 35, and w(5;5;2) = 178. A list of other known exact values of w(k;m;2) appears in [15]. Improved lower bounds on several specific values of w(k;k;s) are given in [3] and [10].

In another direction, Graham [7] gives an elegant proof that if one defines $w_1(k;3)$ to be the least n such that every 2-coloring of [1;n] gives either k consecutive integers in the first color or a 3-term arithmetic progression in the second color, then

$$k^{c\log k} < W_1(k;3) < k^{dk^2};$$

for suitable constants c; d > 0. This immediately gives $w(k/3; 2 < k^{dk^2}$ since we trivially have w(k/3; 2) $w_1(k/3)$. In view of Graham's bounds on $w_1(k/3)$, it would be desirable to obtain improved bounds on w(k/3; 2). Of particular interest is the question of whether or not there is a non-polynomial lower bound for w(k/3; 2).

In this note we give a lower bound of $w(k;3;2) > k^{(2-o(1))}$. Although this may seem weak, we do know that $w(k;3;2) < k^2$ for 5 k 16 (i.e., for all known values of w(k;3;2) with k 5; see Table 2). More generally, we give a lower bound on w(k;m;2) for arbitrary fixed m. We also present a lower bound for the classical van der Waerden numbers w(k;k;:::;k;s) that is a slight improvement over previously published bounds. In addition, we present an upper bound for w(k;4;2) and an upper bound for w(4;4;:::;4;s).

2 Upper and lower bounds for certain van der Waerden functions

We shall need several definitions, which we collect here.

For positive integers k and n,

$$r_k(n) = \max_{S \ [1,n]} f_j S_j$$
: S contains no k-term arithmetic progression g:

For positive integers k and m, denote by $c_k(m)$ the minimum number of colors required to color [1; m] so that there is no monochromatic k-term arithmetic progression.

The function $w_1(k;3)$ has been defined in Section 1. Similarly, we define $w_1(k;4)$ to be the least *n* such that every 2-coloring of [1;n] has either *k* consecutive integers in the first color or a 4-term arithmetic progression in the second color.

We begin with an upper bound for $w_1(k; 4)$. The proof is essentially the same as the proof given by Graham [7] of an upper bound for $w_1(k; 3)$. For completeness, we include the proof here. We will make use of a recent result of Green and Tao [9], who showed that for some constant c > 0,

$$r_4(n) < ne^{-c^{D_{\log\log n}}}$$
(1)

for all *n* 3.

Proposition 1. There exists a constant c > 0 such that $w_1(k;4) < e^{k^{clogk}}$ for all k = 2.

Proof. Suppose we have a 2-coloring of [1; n] (assume n = 4) with no 4-term arithmetic progression of

the second color and no *k* consecutive integers of the first color. Let $t_1 < t_2 < \cdots < t_m$ be the integers of the second color. Hence, $m < r_4(n)$. Let us define $t_0 = 0$ and $t_{m+1} = n$. Then there must be some *i*, 1 = i = m, such that

$$t_{i+1} \quad t_i > \frac{n}{2r_4(n)}$$

(Otherwise, using $r_4(n)$ 3, we would have $n = \sum_{i=0}^{m} (t_{i+1} \quad t_i) \quad \frac{n(m+1)}{2r_4(n)} \quad \frac{n(r_4(n)+1)}{2r_4(n)} \quad \frac{n}{2} + \frac{n}{6}$.) Using (1), we now have an *i* with

$$t_{i+1}$$
 $t_i > \frac{n}{2r_4(n)} > \frac{1}{2}e^{c^{D_{\overline{\log\log n}}}}$

If $n = e^{k^{d \log k}}$, $d = c^2$, then $\frac{1}{2}e^{c^{p} \log \log n}$ k and we have k consecutive integers of the first color, a contradiction. Hence, $n < e^{k^{d \log k}}$ and we are done.

Clearly $w(k;4;2) = w_1(k;4)$. Consequently, we have the following result.

Corollary 2. There exists a constant d > 0 such that $w(k;4;2) < e^{k^{d\log k}}$ for all k = 2.

Using Green and Tao's result, it is not difficult to obtain an upper bound for $w(4;4;\ldots;4;s)$.

Proposition 3. There exists a constant d > 0 such that $w(4;4;\ldots;4;s) < e^{s^{d\log s}}$ for all s = 2.

Proof. Consider a $c_4(m)$ -coloring of [1/m] for which there is no monochromatic 4-term arithmetic progression. Some color must be used at least $\frac{m}{c_4(m)}$ times, and hence $\frac{m}{c_4(m)} = r_4(m)$ so that $\frac{m}{r_4(m)} = c_4(m)$. Let c > 0 be such that (1) holds for all n = 3, and let $m = e^{s^{dlogs}}$, where $d = c^{-2}$. Then $c_4(m) = \frac{m}{r_4(m)} > e^{c^{P} \log \log m} = s$. This means that every s-coloring of [1/m] has a monochromatic 4-term arithmetic progression. Since $m = e^{s^{dlogs}}$, the proof is complete.

It is interesting that the bounds in Corollary 2 and Proposition 3 have the same form.

The following theorem gives a lower bound on w(k;k;...;k;s). It is deduced without too much difficulty from the Symmetric Hypergraph Theorem as it appears in [8], combined with an old result of Rankin [17]. To the best of our knowledgg

Proof. We make use of the observation that for positive integers *s* and *m*, if $s = c_k(m)$, then w(k;k;...;k;s) > m, which is clear from the definitions. For large enough *m*, (2) gives

$$c_k(m) < \frac{2m\log m}{r_k(m)} + \frac{1}{2} = \frac{3m\log m}{r_k(m)}$$
 (4)

Now let $d = \frac{1}{2c} z^{z+1}$, where *c* is from (3), and let $m = s^{d(\log s)^z}$, where *s* is large enough so that (4) holds. By (3), noting that $\log m = d(\log s)^{z+1} = \frac{\log s}{2c} z^{z+1}$, we have

$$\frac{m}{r_k(m)} < e^{c(\log m)^{\frac{1}{2+1}}} = e^{c\frac{\log s}{2c}} = P_{\overline{s}}.$$

Therefore,

$$\frac{3m\log m}{r_k(m)} < 3d'^{D}\bar{s}(\log s)^{z+1} < s$$

for sufficiently large s. Thus, for sufficiently large s,

$$c_k(m) < \frac{3m\log m}{r_k(m)} < s.$$

According to the observation at the beginning of the proof, this implies that $W(k;k;:::;k;s) > m = s^{d(\log s)^{Z}}$, as required.

We now give a lower bound on w(k; m; 2). We make use of the Lovász Local Lemma (see [8] for a proof), which will be implicitly stated in the proof.

Theorem 5. Let *m* 3 be fixed. Then for all sufficiently large *k*,

$$W(k;m;2) > k^{m-1} \frac{1}{\log \log k}$$

Proof. Given *m*, choose k > m large enough so that

$$k^{\frac{1}{2m\log\log k}} > m \quad \frac{1}{2\log\log k} \quad \log k \tag{5}$$

and

$$6 < \frac{\log k}{\log \log k}$$
 (6)

Next, let $n = k^{m-1} k^{k-1}$. To prove the theorem, we will show that there exists a (red, blue)-coloring of [1,n] for which there is no red *k*-term arithmetic progression and no blue *m*-term arithmetic progression.

For the purpose of using the Lovász Local Lemma, randomly color [1;n] in the following way. For each *i* 2 [1;n], color *i* red with probability $p = 1 - k^{a-1}$ where

$$a := \frac{1}{2m \log \log k}$$

and color it blue with probability 1 *p*.

Let P be any k-term arithmetic progression. Then, since $1 + x = e^x$, the probability that P is red is

$$p^{k} = (1 \quad k^{a-1})^{k} \quad e^{k^{a-1} \quad k} = e^{k^{a}}:$$

Hence, applying (5), we have

Also, for any *m*-term arithmetic progression \mathcal{Q} , the probability that \mathcal{Q} is blue is

$$(1 \quad p)^m = (k^{a-1})^m = \frac{1}{k^{m-\frac{1}{2\log\log k}}}$$

Now let P_1 ; P_2 ;...; P_t be all of the arithmetic progressions in [1;n] with length k or m. So that we may apply the Lovász Local Lemma, we form the "dependency graph" G by setting $V(G) = fP_1; P_2; ...; P_tg$ and $E(G) = ffP_i; P_jg: i = j; P_i \setminus P_j \notin \otimes g$. For each $P_i \ge V(G)$, let $d(P_i)$ denote the degree of the vertex P_i in G, i.e., *jfe* $\ge E(G) : P_i \ge egj$. We now estimate $d(P_i)$ from above. Let $x \ge [1;n]$. The number of k-term arithmetic progressions P in [1;n] that contain x is bounded above by $k = \frac{n}{k-1}$. Similarly, the number of m-term arithmetic progressions Q in [1;n] that contain x is bounded above by $m = \frac{n}{m-1}$.

Let P_i be any *k*-term arithmetic progression contained in [1;n]. The total number of *k*-term arithmetic progressions P and *m*-term arithmetic progressions O in [1;n] that may have nonempty intersection with P_i is bounded above by

$$k \quad k \quad \frac{n}{k-1} + m \quad \frac{n}{m-1} \quad < kn \quad 2 + \frac{2}{m-1} \quad ;$$
 (7)

since k > m. Thus, $d(P_i) < kn(2 + \frac{2}{m-1})$ when $jP_ij = k$. Likewise, $d(P_i) < mn(2 + \frac{2}{m-1})$ when $jP_ij = m$. Thus, for all vertices P_i of G, we have $d(P_i) < kn(2 + \frac{2}{m-1})$.

To finish setting up the hypotheses for the Lov'asz Local Lemma, we let X_i denote the event that the arithmetic progression P_i is red if $jP_i j = k$, or blue if $jP_i j = m$, and we let $d = \max_1 t d(P_i)$. We have seen above that for all i, 1 = i t, the probability that X_i occurs is at most

$$q := \frac{1}{k^{m} \frac{1}{2\log\log k}};$$

while from (7) we have $d < 2kn(1 + \frac{1}{m-1})$.

We are nowww.1

Table 1: Small values of W(k;3) and $W_1(k;3)$

Table 1. Small values of $W(x, 3)$ and $W_1(x, 3)$															
k	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
W(k;3;2)	6	9	18	22	32	46	58	77	97	114	135	160	186	218	238
$W_1(k;3)$	9	23	34	73	113	193	?	?	?	?	?	?	?	?	?

d < 3kn, so that d+1 < 3kn+

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